

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: 1431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

**IN THE CLAIMS**

Please cancel claims 1-29, 54, and 61-63.

Please amend claims 30, 32-35, 37-48, 51-53, and 55-60 as follows:

1-29. (Cancelled)

30. (Currently Amended) An evaluation circuit comprising:

a first linear automation circuit;

a second linear automation circuit connected in parallel with the first linear automation circuit, each having a set of states, which have a common input line for receiving a data stream,  $T_n$  comprising  $n$  successive data words  $y(1), \dots, y(n)$ , each having a width of  $k$  bits, wherein the first linear automaton circuit and the second linear automation circuit are configured such that a first signature can be calculated from each data word and a second signature can be calculated from each data word;

a first logic combination gate and a second logic combination gate that compare the first signature and the second signature, respectively, with a predeterminable good signature and an output comparison value.

31. (Previously Presented) The evaluation circuit as claimed in claim 30, comprising wherein the first logic combination gate and the second combination logic gate are exclusive-OR gates having first inputs, respectively, connected to the outputs of the associated first and second linear automaton circuit and to whose second inputs good signatures can be applied.

32. (Currently Amended) The evaluation circuit as claimed in claim 30, comprising wherein arranged upstream of the first linear automaton circuit is a first coder, that ~~codes~~ encodes ~~the~~ a data word having the data word length of  $k$  bits into an ~~coded~~ encoded data word  $u^l(i)$ ,

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: 1431-135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

$u^1(i) = \text{Cod1}$  having the word width of  $K1$  bits, for  $i=1, \dots, n$ , and where  $\text{Cod1}$  represents the ~~encoding~~ encoding function of the first coder.

33. (Currently Amended) The evaluation circuit as claimed in claim 32, comprising wherein the following holds true for the ~~encoding~~ encoding function of the first coder (C1):

$$\text{Cod1}(y'(i)) = u^1(i) \oplus f_1(e(i)),$$

or

$$\text{Cod1}(y'(i)) = \text{Cod1}(y(i) \oplus e(i)) = \text{Cod1}(y(i) \oplus f_1(e(i)))$$

where a function  $f_1$  by  $f_1(0) = 0$  exists for  $y'(i) = y(i) \oplus e(i)$ , and where a function  $f_1^{-1}$  where

$$f_1^{-1}(f_1(e)) = e$$

exists for all binary data words  $e$  having the word width  $k$  which may occur as errors of a data word, where  $e$  denotes a faulty data word of the data stream  $T_n$ .

34. (Currently Amended) The evaluation circuit as claimed in ~~one of~~ claim 30, comprising wherein arranged upstream of the second linear automaton circuit is a second coder, which ~~codes~~ encodes the data word  $y(i)$  having the data word length of  $k$  bits into a ~~coded~~ encoded data word  $u^2(i)$ ,  $u^2(i) = \text{Cod2}(y(i))$  having the word width of  $K2$  bits, for  $i=1, \dots, n$ , and where  $\text{Cod2}$  represents the ~~encoding~~ encoding function of the second coder (C2).

35. (Currently Amended) An evaluation circuit for detecting and/or locating faulty data words in a data stream  $T_n$  comprising:

a first linear automaton circuit and a second linear automaton circuit connected in parallel, each having a set of states, wherein the first linear automaton circuit and the second linear automaton circuit have a common input line for receiving a data stream  $T_n$  comprising  $n$  successive data words  $y(1), \dots, y(n)$  each having a width of  $k$  bits,

wherein the first linear automaton circuit can be described by the following equation

$$z(t+1) = Az(t) \oplus y(t)$$

wherein the second linear automaton circuit can be described by the following equation

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

$$z(t+1) = Bz(t) \oplus y(t)$$

where A and B represent the state matrices of the linear automaton circuits, where the state matrices A and B can be inverted, and where the dimension L of the state vectors is  $\geq k$ ,

the first linear automaton circuit and the second linear automaton circuit are designed such that a first signature and a second signature, respectively, can be calculated of each data word,

L first logic combination gates arranged downstream of the first linear automaton circuit and also L second logic combination gates arranged downstream of the second linear automaton circuit,

the logic combination gates are designed such that the signature respectively calculated by the linear automaton circuit can be compared with a predeterminable good signature and a comparison value can be output.

36. (Previously Presented) The evaluation circuit as claimed in claim 35, comprising wherein the logic combination gates are present as exclusive-OR gates whose first inputs are respectively connected to the outputs of the associated linear automaton circuit (L1, L2) and to whose second inputs good signatures can be applied.

37. (Currently Amended) The evaluation circuit as claimed in claim 35, comprising wherein arranged upstream of the first linear automaton circuit is a first coder, that ~~codes~~ encodes the data word  $y(i)$  having the data word length of k bits into a ~~coded~~ an encoded data word  $u^1(i)$ ,  $u^1(i) = \text{Cod1}$  having the word width of  $K1$  bits, for  $i=1, \dots, n$ , and where Cod1 represents the ~~coding~~ encoding function of the first coder.

38. (Currently Amended) The evaluation circuit as claimed in claim 37, comprising wherein the following holds true for the ~~coding~~ encoding function of the first coder:

$$\text{Cod1}(y'(i)) = u^1(i) \oplus f_1(e(i)),$$

or

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

$$\text{Cod1}(y'(i)) = \text{Cod1}(y(i) \oplus e(i)) = \text{Cod1}(y(i) \oplus f_1(e(i)))$$

where a function  $f_1$  by  $f_1(0) = 0$  exists for  $y'(i) = y(i) \oplus e(i)$ , and where a function  $f_1^{-1}$  where

$$f_1^{-1}(f_1(e)) = e$$

exists for all binary data words  $e$  having the word width  $k$  which may occur as errors of a data word, where  $e$  denotes a faulty data word of the data stream  $T_n$ .

39. (Currently Amended) The evaluation circuit as claimed in ~~one of claim 35~~ claim 37, comprising wherein arranged upstream of the second linear automaton circuit is a second coder, which encodes the data word  $y(i)$  having the data word length of  $k$  bits into an encoded data word  $u^2(i)$ ,  $u^2(i) = \text{Cod2}(y(i))$  having the word width of  $K2$  bits, for  $i=1, \dots, n$ , and where  $\text{Cod2}$  represents the encoding function of the second coder.

40. (Currently Amended) The evaluation circuit as claimed in claim 39, comprising wherein the following holds true for the encoding function of the second coder:

$$\text{Cod2}(y'(i)) = u^2(i) \oplus f_2(e(i)),$$

or

$$\begin{aligned} \text{Cod2}(y'(i)) &= \text{Cod2}(y(i) \oplus e(i)) \\ &= \text{Cod2}(y(i)) \oplus f_2(e(i)) \end{aligned}$$

where a function  $f_2^{-1}$  where

$$f_2^{-1}(f_2(e)) = e$$

exists for all binary data words  $e$  having the word width  $k$  which may occur as errors of a data word, where  $e$  denotes a faulty data word of the data stream  $T_n$ .

41. (Currently Amended) The evaluation circuit as claimed in ~~one of claim 37~~ claim 39, comprising wherein that the word width  $K1$  of the data words  $u^1(i)$  encoded by the first coder is equal to the word width  $K2$  of the data words  $u^2(i)$  encoded by the second coder.

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

42. (Currently Amended) The evaluation circuit as claimed in ~~one of claim 3739~~, comprising wherein

the first coder matches the second coder with regard to its construction and its function.

43. (Currently Amended) The evaluation circuit as claimed in ~~one of claim 3739~~, comprising wherein

the word width  $K1$  of the data words  $u^1(i)$  encoded by the first coder and the word width  $K2$  of the data words  $u^2(i)$  encoded by the second coder are in each case equal to the word width  $k$  of the data words  $y(1), \dots, y(n)$  of the data stream  $T_n$ .

44. (Currently Amended) The evaluation circuit as claimed in ~~one of claim 3739~~, comprising wherein

the encoding functions  $Cod1$  and  $Cod2$  of the first coder and of the second coder are designed as follows:

$$\begin{aligned} &Cod1(y_1(i), y_2(i), \dots, y_k(i)) \\ &= P1(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0) \end{aligned}$$

$$\begin{aligned} &Cod2(y_1(i), y_2(i), \dots, y_k(i)) \\ &= P2(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0) \end{aligned}$$

for  $i, 1, \dots, n$

where the number of zeros situated at the end of  $P1(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0)$  is equal to  $(K1-k)$ , where the number at the end of  $P2(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0)$  is equal to  $(K2-k)$ , and where  $P1$  represents an arbitrary permutation of the  $K1$  components of  $(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0)$  and  $P2$  represents an arbitrary permutation of the  $K2$  components of  $(y_1(i), y_2(i), \dots, y_k(i), 0, \dots, 0)$ .

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: 1431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

45. (Currently Amended) The evaluation circuit as claimed in claim 37, comprising wherein the coding functions Cod1 and Cod2 of the first coder and of the second ~~encoder~~ coder are designed as follows:

$$\begin{aligned} &\text{Cod1}(y_1(i), y_2(i), \dots, y_k(i)) \\ &= P1(y_1(i), y_2(i), \dots, y_k(i), b_1^1 \dots, b_{K1-k}^1) \end{aligned}$$

$$\begin{aligned} &\text{Cod2}(y_1(i), y_2(i), \dots, y_k(i)) \\ &= P2(y_1(i), y_2(i), \dots, y_k(i), b_1^2 \dots, b_{K2-k}^2) \end{aligned}$$

where  $b_1^1, \dots, b_{K1-k}^1, b_1^2, \dots, b_{K2-k}^2 \in \{0,1\}$ , and where P1 and P2 represent arbitrary permutations.

46. (Currently Amended) The evaluation circuit as claimed in ~~one of~~ claim 37, comprising wherein the encoding function Cod1 of the first coder is designed such that it realizes a linear block code,  $f_1 = \text{Cod1}$ .

47. (Currently Amended) The evaluation circuit as claimed in ~~one of~~ claim ~~37~~39, comprising wherein the encoding function Cod2 of the second coder is designed such that it realizes a linear block code,  $f_2 = \text{Cod2}$ .

48. (Currently Amended) The evaluation circuit as claimed in ~~one of~~ claim 35, comprising wherein the state matrix A of the first linear automaton circuit and the state matrix B of the second linear automaton circuit are related to one another as follows:

$$B = A^n$$

where  $n \neq 1$ .

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

49. (Previously Presented) The evaluation circuit as claimed in claim 35, comprising wherein the state matrix B of the second linear automaton circuit is equal to the inverted state matrix  $A^{-1}$  of the first linear automaton circuit.

50. (Previously Presented) The evaluation circuit as claimed in claim 35, comprising wherein the first linear automaton circuit is designed as a linear feedback shift register and the second linear automaton circuit is designed as an inverse linear feedback shift register, both linear automaton circuits having a parallel input.

51. (Currently Amended) The evaluation circuit as claimed in claim 35, comprising wherein the first linear automaton circuit is designed as a linear feedback, ~~K1-dimensional~~ multi-input shift register and/or the second linear automaton circuit is designed as a linear feedback, ~~K2-dimensional~~ multi-input shift register.

52. (Currently Amended) The evaluation circuit as claimed in claim 51, comprising wherein the multi-input shift register/~~registers~~ has/~~have~~ a primitive feedback polynomial of maximum length.

53. (Currently Amended) A method for testing an integrated circuit device by detecting and/or locating faulty data words in a data stream  $T_n$ , the method having the following method steps of:

inputting data words  $y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n)$  of a data stream  $T_n$  into a first coder,

encoding the data words  $y(1), \dots, y(n)$  into encoded data words  $u^1(1), \dots, u^1(n)$  having the a word width K1 where  $K1 \geq k$  by means of the coding function Cod1 of ~~the~~ a first coder,

inputting the coded data words  $u^1(1), \dots, u^1(i-1), u^{1'}(i)$  or  $u^1(i), u^1(i), \dots, u^1(n)$  into the inputs of a first linear automaton circuit, which is described by the automaton equation;



**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_n$ 

$$z^1(t+1)=A \cdot z^1(t)+u^1(t)$$

where  $t$  is an instant in time,  $z^1$  represents a  $K1$ -dimensional state vector and  $A$  represents a  $K1 \times K1$  state matrix, and where the state matrix  $A$  can be inverted,

processing the coded data words  $u^1(1), \dots, u^1(i-1), u^{1'}(i)$  or  $u^1(i), u^1(i), \dots, u^1(n)$  by means of the first linear automaton circuit, the first linear automaton circuit,

undergoing transition to the state  $z^1(n+1)=S_1(L1, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$  if no error can be detected in the case of the coded data words  $u^1(1), \dots, u^1(i-1), u^1(i), u^1(i+1), \dots, u^1(n)$ ,

undergoing transition to the state  $z^{1'}(n+1)=S_1(L1, y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n))$  if an error is present at least in the case of the  $i$ -th position of the coded data words  $u^1(1), \dots, u^1(i-1), u^{1'}(i), \dots, u^1(n)$ ,

the signature of an error-free data stream  $T_n$  being designated by  $S(L1, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$  and the signature of a faulty data stream  $T_n$  being designated by  $S(L1, y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n))$ ,

checking the determined signature of the data stream  $T_n$  and continuing with method step a) for further data streams  $T_n$  if the determined signature of the data stream  $T_n$  is the signature of an error-free data stream  $T_n$ ,

inputting the data words  $y(1), \dots, y(i-1), y'(i), \dots, y(n)$  of the data stream  $T_n$  in a second coder,

coding the data words  $y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n)$  to coded data words  $u^2(1), \dots, u^2(i-1), u^{2'}(i)$  or  $u^2(i), u^2(i), \dots, u^2(n)$  having the word width  $K2$  where  $K2 \geq k$  by means of the coding function  $Cod2$  of the second coder,

inputting the coded data words  $u^2(1), \dots, u^2(i-1), u^{2'}(i)$  or  $u^2(i), u^2(i), \dots, u^2(n)$  into the inputs of a second linear automaton circuit, which is described by the automaton equation

$$z^2(t+1) = B \cdot z^2(t) \oplus u^2(t)$$

where  $z^2$  represents a  $K2$ -dimensional state vector and  $B$  represents a  $K2 \times K2$  state matrix where  $B \neq A$ , and where the state matrix  $B$  can be inverted,



**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_n$ 

---

processing the coded data words  $u^2(1), \dots, u^2(i-1), u^{2'}(i)$  or  $u^2(i), u^2(i), \dots, u^2(n)$  by means of the second linear automaton circuit, the second linear automaton circuit,

undergoing transition to the state  $z^2(n+1)=S_2(L2, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$  if no error can be detected in the case of the data words  $u^2(1), \dots, u^2(i-1), u^2(i), u^2(i), \dots, u^2(n)$ ,

undergoing transition to the state  $z^{2'}(n+1)=S_2(L2, y(1), \dots, y(i-1), y(i), y'(i), y(i+1), \dots, y(n))$  if an error is present at least in the case of the  $i$ -th position of the coded data words  $u^2(1), \dots, u^2(i-1), u^{2'}(i), u^2(i), \dots, u^2(n)$ ,

the signature of an error-free data stream  $T_n$  being designated by  $S(L2, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$  and the signature of a faulty data stream  $T_n$  being designated by  $S(L2, y(1), \dots, y(i-1), y'(i), \dots, y(n))$ ,

determining the signature differences  $\Delta S1$  and  $\Delta S2$  by means of exclusive-OR logic combinations of the signatures  $S1$  and  $S2$  ~~determined in method step d) and i), respectively,~~ with ascertained good signatures, in each case according to the following specifications:

$$\Delta S1 = S(L1, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$$

$$\oplus S(L1, y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n))$$

$$\Delta S2 = S(L2, y(1), \dots, y(i-1), y(i), y(i+1), \dots, y(n))$$

$$\oplus S(L2, y(1), \dots, y(i-1), y'(i), y(i+1), \dots, y(n))$$

determining a unique solution for the position  $i$  of the faulty bit in the faulty data word by solving the equation

$$f_1^{-1}(A^{i-n}\Delta S1) = f_2^{-1}(B^{i-n}\Delta S2)$$

and if no unique solution results for  $1 \leq i \leq n$ , outputting a notification by means of an output medium that two or more errors are present in the data stream  $T_n$  under consideration,

determining a unique solution for the counter  $e(i)$  of the faulty data word  $y'(i)$  in the data stream  $T_n$  by solving the equation

$$e(i) = f_1^{-1}(A^{i-n} \cdot \Delta S1)$$

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: I431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

outputting the position  $i$  of the faulty bit in the faulty data word and also the error  $e(i)$  of the faulty data word  $y'(i)$  in the data stream  $T_n$  by means of an output medium; and  
evaluating an integrated circuit in response to the output.

54. (Cancelled)

55. (Currently Amended) The evaluation circuit as claimed in ~~one of~~ claim 35, comprising wherein the evaluation circuit is monolithically integrated on an integrated circuit, each data word having a data word length of  $k$ .

56. (Currently Amended) A loadboard for receiving at least one needle card for testing integrated circuits ~~and/or~~ having at least one test socket for testing integrated circuits ~~and/or~~ for connecting a handler to a tester of integrated circuits, the loadboard having an evaluation circuit as claimed in ~~one of~~ claim 35.

57. (Currently Amended) A needle card for testing integrated circuits, in which an evaluation circuit as claimed in ~~one of~~ claim 35 is integrated.

58. (Currently Amended) A tester for testing integrated circuits having the following features:

the tester is provided with a plurality of instruments for generating signals or data streams and with a plurality of measuring sensors, in particular for currents and voltages,

the tester has a loadboard which is provided for receiving at least one needle card for testing integrated circuits ~~and/or~~ for connecting a handler to a tester of integrated circuits ~~and/or~~ which is equipped with at least one test socket for testing integrated circuits, and

the tester has an evaluation circuit as claimed in ~~one of~~ claim 35.

**Amendment and Response**

Applicant: Michael Goessel et al.

Serial No.: 10/577,288

Filed: April 24, 2006

Docket No.: 1431.135.101/FIN516PCT/US

Title: EVALUATION CIRCUIT AND METHOD FOR DETECTING AND/OR LOCATING FAULTY DATA WORDS IN A DATA STREAM  $T_N$

---

59. (Currently Amended) A storage medium storing a computer program for executing a method for detecting and/or locating faulty data words in a data stream  $T_n$ , which is designed for executing a portion of the method of claim 53.

60. (Currently Amended) The storage medium computer program as claimed in claim 59, wherein the which is contained on a storage medium, in particular in a is a computer memory or in a random access memory.

61-63 (Cancelled)